

Endpoint-Explicit Differential Dynamic Programming via Exact Resolution

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Various approaches have emerged in recent years to address the challenge of parallelizing optimal control (OC) algorithms. To achieve this, it is critical to handle endpoint constraints efficiently. The **contribution of our work** are the follows

- A novel endpoint-explicit DDP, which is both exact and computationally efficient, offering quadratic convergence.
- Our method **effectively handles rank deficiencies** without increasing computational cost.

(1)

• Highly suitable for **MPC applications** with both **forward and inverse dynamics** formulations.

OC problem abstraction

OC problems subject to stagewise and endpoint constraints can be locally resolved by solving the

Endpoint-explicit DDP

Nonlinear OC problems with endpoint constraints can be formulated using either forward or inverse dynamics:

following quadratic program:

$$\min_{\mathbf{w}} \frac{1}{2} \mathbf{w}^{\mathsf{T}} \mathbf{A} \mathbf{w} - \mathbf{w}^{\mathsf{T}} \mathbf{a}$$

subject to $\mathbf{B} \mathbf{w} = \mathbf{b}$.

In this abstraction, $\mathbf{w} \in \mathbb{R}^{n_w}$ represents the primal and dual decision variables encountered in an OC problem without endpoint constraints (e.g., search direction for states, controls, and multipliers of the dynamics constraints). Moreover, $\mathbf{B} \in \mathbb{R}^{n_b \times n_w}$ represents the Jacobian of the endpoint constraint, while $\mathbf{A} \in S_{++}^{n_w}$ is a large banded matrix that describes the OC problem without endpoint constraints.

The solution to this system is:

$$\mathbf{y} = -\mathbf{S}^{-1}(\mathbf{b} - \mathbf{B}\mathbf{A}^{-1}\mathbf{a}), \qquad (2a)$$
$$\mathbf{w} = \mathbf{A}^{-1}\mathbf{a} - \mathbf{A}^{-1}\mathbf{B}^{\mathsf{T}}\mathbf{y}, \qquad (2b)$$

N - 1 $\min_{\mathbf{x}_s,\mathbf{u}_s} \ell_N(\mathbf{x}_N) + \sum_{k=0}^{N-1} \ell_k(\mathbf{x}_k,\mathbf{u}_k)$ s.t. $\tilde{\mathbf{x}}_0 \ominus \mathbf{x}_0 = \mathbf{0}$, (initial condition) $\mathbf{f}_k(\mathbf{x}_k, \mathbf{u}_k) \ominus \mathbf{x}_{k+1} = \mathbf{0}$, (integrator / forward dyn.) $\mathbf{h}_k(\mathbf{x}_k, \mathbf{u}_k) = \mathbf{0}$, (inverse dyn.) $\mathbf{r}(\mathbf{x}_N) = \mathbf{0}$. (endpoint constraint) (3)In general, our algorithm can be divided into three main steps: 1. We compute the search direction for the **endpoint-independent** problem $= A^{-1}a$) through Riccati recursion, performing the unconstrained DDP algorithm: $(\hat{\mathbf{w}})$

$$\begin{split} \delta \hat{\mathbf{u}} &= -\hat{\boldsymbol{\pi}} - \boldsymbol{\Pi} \delta \hat{\mathbf{x}}, \\ \delta \hat{\mathbf{x}}_{k+1} &= \mathbf{f}_{\mathbf{x}_k} \delta \hat{\mathbf{x}}_k + \mathbf{f}_{\mathbf{u}_k} \delta \hat{\mathbf{u}}_k + \overline{\mathbf{f}}_{k+1}, \quad \delta \hat{\mathbf{x}}_0 &= \overline{\mathbf{f}}_0, \end{split} \begin{aligned} \hat{\boldsymbol{\pi}}_u &\coloneqq \mathbf{Q}_{\mathbf{u}\mathbf{u}}^{-1} \mathbf{Q}_{\mathbf{u}\mathbf{u}} = \mathbf{Q}_{\mathbf{u}\mathbf{u}}^{-1} \mathbf{Q}_{\mathbf{u}\mathbf{u}} &\text{(forward dyn.)} \\ \hat{\boldsymbol{\pi}}_s &\coloneqq \mathbf{k} + (\mathbf{k}_s^{\top} \mathbf{u}_{\mathbf{u}} \Psi_s^{\top})^{\top}, \ \boldsymbol{\Pi}_s &\coloneqq \mathbf{K} + (\mathbf{K}_s^{\top} \mathbf{u}_{\mathbf{u}} \Psi_s^{\top})^{\top} \\ &(\text{inverse dyn.}) \end{aligned}$$

2. We run another Riccati recursion to compute the endpoint-dependent search direction ($\dot{W} =$ $A^{-1}B^{\intercal}$). This step reuses the factorization computed in (1.), and replaces a by B^{\intercal} in Equation (1):

$$\delta \check{\mathbf{U}}_{\mathbf{c}} = -\check{\boldsymbol{\pi}}_{\mathbf{c}} - \boldsymbol{\Pi} \delta \check{\mathbf{X}}_{\mathbf{c}}, \qquad \qquad \check{\boldsymbol{\pi}}_{\mathbf{c}_{s}} \coloneqq \mathbf{Q}_{\mathbf{u}\mathbf{u}} \mathbf{Q}_{\mathbf{u}\mathbf{c}} \text{ (forward dyn.)} \\ \check{\boldsymbol{\pi}}_{\mathbf{c}_{s}} \coloneqq \mathbf{k}_{\mathbf{c}} + (\mathbf{k}_{s\mathbf{c}}^{\mathsf{T}} \tilde{\mathbf{Q}}_{\mathbf{u}\mathbf{u}} \Psi_{s}^{\mathsf{T}})^{\mathsf{T}} \text{ (inverse dyn.)} \\ \delta \check{\mathbf{X}}' - \mathbf{f} \delta \check{\mathbf{X}} + \mathbf{f} \delta \check{\mathbf{U}} \quad \delta \check{\mathbf{X}} = \mathbf{0} \qquad (5)$$

where $\mathbf{S} = \mathbf{B}\mathbf{A}^{-1}\mathbf{B}^{\mathsf{T}}$ denotes the Schur complement, and its inversion can be efficiently computed via a sparse Cholesky decomposition. Here, Equation (2a) calculates the Lagrange multiplier associated with the endpoint constraint. In contrast, Equation (2b) computes the **primal and dual variables**. This is achieved through two **Riccati recursions** with different initializations. This abstract solution results in a highly efficient algorithm capable of handling arbitrary stagewise constraints.



 $A^{-1}B^{\mathsf{T}}y$

$$\delta \mathbf{u} = \delta \hat{\mathbf{u}} - \delta \check{\mathbf{U}}_{\mathbf{c}} \boldsymbol{\beta}^{+},$$
$$\delta \mathbf{x}' = \delta \hat{\mathbf{x}}' - \delta \check{\mathbf{X}}'_{\mathbf{c}} \boldsymbol{\beta}^{+}.$$

Results: Cost and feasibility, Endpoint factorization, Gymnastic Maneuvers and MPC trials

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Costs and endpoint feasibility We compared **forward** and **inverse dynam**ics formulations to evaluate our endpointexplicit strategy on problems with different stagewise constraints. Here, inverse dynamics generally converges faster, while **cost** and feasibility trends vary by problem.

Compared to the Schur complement method, it achieves similar computation times using LU or QR with pivoting.



MPC trials

To demonstrate its relevance for control tasks, we experimentally validated our MPC controller using our endpoint-explicit method on scenarios like: a B1quadruped with a Z1 manipulator.





Gymnastic maneuvers

Also, our endpoint-explicit DDP algorithm efficiently solved complex gymnastic maneuvers for the Talos humanoid, achieving faster convergence, shorter runtimes, and better constraint satisfaction.







We compared tracking error between the traditional cost-based penalty method and our approach with an explicit terminal constraint. In both the simulated and real-world scenarios, the **terminal constraint** consistently improved tracking performance.

Endpoint factorization

the other hand, we evaluated On our nullspace factorization's ability to handle rank-deficient constraints efficiently.

